TITLE OF THE INVENTION

STORAGE SYSTEM

BACKGROUND OF THE INVENTION

Field of the Invention

The present invention relates to a storage system. More particularly, the present invention relates to a method of communicating data between storage apparatus such as RAIDs connected to each other by a network or between a storage apparatus and a host, which are also connected to each other by a network, and relates to the configuration of a storage device for implementing the method.

Description of the Related Art

Popularization of IP storages and a rising demand for remote copying of data have resulted in more frequent data transfers between storage apparatus, which are separated from each other by long distances and connected to each other by a network, or between a storage apparatus and a server, which are also separated from each other by long distances and connected to each other by a network. In addition, the amount of transferred data has been increasing year after year so that it becomes necessary to transmit data having a large amount at a high speed with a high degree of reliability.

As an interface between a server and a storage apparatus connected to the server, a Fibre Channel and an SCSI interface are known. The Fibre Channel is adopted as a standard interface, which composes a SAN (Storage Area Network) and is used for transferring data at a high speed. On the other hand, the SCSI interface offers merits of a high data transmission speed, a short transmission delay and an extremely low probability of transmission error generation. Thus, the SCSI interface is used as a popular conventional protocol for storage apparatus. However, the SCSI interface has a problem of limiting applications of the SCSI interface to only communications at short transmission distances.

In recent years, trials have been made to connect storage apparatus to each other by adopting an iSCSI protocol. The iSCSI protocol is a protocol technology for implementing SCSI processing as an interface technology on a TCP/IP, which is a network technology. At the present time, the IETF (the Internet Engineering Task Force) is carrying out work of standardization of the iSCSI protocol. It is expected that there will be many cases in which packets are lost at an IP layer when data is transferred by adoption of the iSCSI protocol between storage apparatus separated from each other by a long distance.

In general, as a countermeasure for tackling the

phenomenon of losing packets in communication of data,

there are known an ARQ (Automatic Repeat reQuest) technique,

i.e., an automatic retransmission request technique, and an

FEC (Forward Error Correction) technique. When data is

transmitted over a long distance, it is quite within the

bounds of possibility that packets of the data are lost.

In consequence, when data is communicated by adoption of

the ARQ technique, it takes time to retransmit data and, as

a result, the data transmission efficiency decreases. In

particular, a network delay included in the time it takes

to retransmit data in a data transmission over a long

distance increases in proportion to the distance. For this

reason, the ARQ technique cannot be said to be a desirable

countermeasure.

For example, in Japanese Patent Laid-open No. 2001-

7785, there is disclosed a technology adopting both the FEC

and ARQ techniques. In accordance with this technology, a

communication partner is informed of an error control

technique suitable for the line quality of the transmission

network by adding a unique word to transmitted data and, on

the receiver side, the data is decoded in dependence on the

unique word. In addition, in Japanese Patent Laid-open No.

2001-168944, there has been proposed a technique for

transmitting 2 or more types of data having different

attributes, namely, ISO data and ASY data, and, on the

receiver side, processing is carried out to receive the data in accordance with the attributes.

However, the above publications neither describe a relation with storage apparatus nor disclose ways to keep up with changes in FEC redundancy.

SUMMARY OF THE INVENTION

It is thus an object of the present invention to provide a communication method that is capable of restoring transmitted data in a data transmission through a network even if packets of the data are lost.

It is another object of the present invention to provide a storage apparatus conforming to the iSCSI protocol adopting the FEC technique.

It is a further object of the present invention to provide a storage system that is capable of transmitting data by changing the status of an FEC process and the redundancy of a data transmission between iSCSI layers in accordance with the destination of the data transmission and capable of restoring the transmitted data.

In order to achieve the objects described above, the present invention provides a system with each of storage apparatus thereof adopting an iSCSI protocol and each data exchanged between iSCSI layers in a storage system in which the storage apparatus are connected to each other by a

network or the storage apparatus are connected to a host computer by a network. The host computer has the general meaning of a computer including a server. Each of the storage apparatus has an FEC control unit. On the transmitter side, the FEC control unit carries out redundancy conversion processing on data to be transmitted in an encoding process and transmits a group of redundancy packets of the data completing the redundancy conversion processing to the network. On the other hand, by using the group of redundancy packets received from the network as a base, the receiver restores the received data in FEC units in a decoding process. By carrying out the encoding process to carry out the redundancy conversion processing on the data to be transmitted and the decoding process to restore the received data as described above, the original group of packets can be restored even if some of redundancy packets in the group are lost and, thus, cannot be received. If the data cannot be restored on the receiver side, on the other hand, the transmitter can retransmit the data typically upon detection of a monitored timeout. The detected timeout is caused by the fact that the receiver does not transmit an ACK to the transmitter.

The present invention is capable of changing the FEC redundancy for packets to be transmitted. To be more specific, the state of redundancy on the transmitter side

is changed in accordance with the state of the packet loss on the receiver side. As a typical means for changing the redundancy, for example, it is possible to change the ratio of the amount of data subjected to the encoding process to carry out redundancy processing on data to be transmitted to the amount of data to be exchanged. In addition, a packet loss ratio can be acquired for each destination of transmission. The packet loss ratio is defined as a ratio of the number of lost packets to the number of all packets in the group, which are supposed to be received by the receiver. The transmitter then transmits data to a destination of transmission by changing the redundancy of the data on the basis of the packet loss ratio acquired from the destination. In the case of a high packet loss ratio, for example, data is transmitted by raising the FEC redundancy. In the case of a low packet loss ratio, on the other hand, data is transmitted by lowering the FEC redundancy. The packet loss ratio is expected to change sometimes in dependence on the network. In this case, the redundancy is changed in accordance with the variations in the packet loss ratio. In addition, it is possible to change intervals at which packets are transmitted in accordance with the packet loss ratio. Thus, in a broad sense, the state of data transmission can be changed in accordance with the packet loss ratio.

A preferred embodiment comprises: a transmission management table for cataloging an FEC redundancy for each destination of transmission and for managing the FEC redundancies; a reception management table for cataloging an FEC redundancy for each source of transmission and for managing the FEC redundancies; an encoding unit for carrying out an FEC encoding process on iSCSI-layer packet data, which is generated by a storage apparatus for a destination of transmission, by referencing the transmission management table and providing the data with a redundancy for the destination of transmission; and a decoding unit for carrying out an FEC decoding process on packet data received from a network in order to decode the packet data back into the iSCSI-layer data. This preferred embodiment can be built as an adapter connected to or embedded in a storage apparatus.

In addition, the present invention is capable of implementing a communication method for transmitting data between iSCSI layers by way of a network in a storage system adopting the iSCSI protocol. The communication method has a first communication mode for transmitting and receiving data in a communication mode including an FEC process, a second communication mode for transmitting and receiving data in a TCP/IP communication mode, forming a judgment as to whether or not a partner serving as a data

communication destination has an iSCSI layer on the basis

of the existence of an iSCSI Name, carrying out an FEC

process based upon an FEC redundancy provided for a

communication partner on data to be transmitted and

transmitting the data completing the FEC process to the

partner in the first communication mode to the partner in a

case where an outcome of the judgment indicates that the

partner has an iSCSI layer, and transmitting data to a

communication partner in the second communication mode in a

case where an outcome of the judgment indicates that the

partner does not have an iSCSI layer.

In accordance with a preferred embodiment

implementing a storage system for communicating data

between iSCSI layers, a transmission of an iSCSI login

frame for starting a TCP session is monitored and, in the

case of a session with a communication partner determined

in advance, the redundancy of data is changed in an

increasing direction when the data is transmitted to the

destination address of the communication partner as long as

this session continues to exist. The communication partner

is an iSCSI node such as an iSCSI initiator or an iSCSI

target.

BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a block diagram showing a storage system

implemented by an embodiment of the present invention.

Fig. 2 is a block diagram showing a typical storage apparatus employing an FEC conversion adapter implemented by an embodiment of the present invention.

Fig. 3 is a block diagram showing a storage system implemented by another embodiment of the present invention.

Fig. 4 is a diagram showing the concept of a data transfer according to an embodiment of the present invention.

Fig. 5 is a diagram showing a typical format of a packet used in communications.

Fig. 6 is a diagram showing the format of an FEC header.

Fig. 7 is a diagram showing typical information cataloged in an FEC control table, an FEC transmission management table and an FEC reception management table.

Fig. 8 shows a flowchart referred to in explaining operations of an FEC encoding process carried out in a transmission of data.

Fig. 9 shows a flowchart referred to in explaining operations carried out in an FEC encoding process for insufficient transmitted data.

Fig. 10 shows a flowchart referred to in explaining operations of a process to retransmit an FEC packet.

Fig. 11 shows a flowchart referred to in explaining

operations of an FEC decoding process in a reception of data.

Fig. 12 shows a continuation flowchart referred to in explaining the operations of the FEC decoding process in the reception of data.

Fig. 13 shows a flowchart referred to in explaining operations of a process carried out in the event of an error in a reception of data.

Fig. 14 shows a flowchart referred to in explaining operations of a process to report a packet loss ratio in a reception of data.

Fig. 15 shows a flowchart referred to in explaining operations of a process carried out for a reported packet loss ratio.

Fig. 16 is a block diagram showing a storage system implemented by a further embodiment of the present invention.

DESCRIPTION OF THE PREFFRED ENBODIMENTS

Embodiments of the present invention will be described below by referring to the drawings.

Fig. 1 is a block diagram showing a storage system connected to a network in accordance with an embodiment of the present invention. As shown in the figure, an IP network 400 is connected to storage apparatus 100 and 200

as well as a host computer 300, which is referred to hereafter simply as the host 300. In this storage system, data having a packet format is exchanged by way of the IP network 400 between the storage apparatus 100 and 200 or between the storage apparatus 100 or 200 and the host 300. Data is communicated between the storage apparatus 100 and 200 typically in a remote copy operation. On the other hand, data is communicated between the storage apparatus 100 or 200 and the host 300 for example in data processing or when the storage apparatus 100 or 200 is used as a data station. The storage system is characterized in that the storage apparatus 100 and 200 and the host 300 are provided with FEC conversion adapters 110, 210 and 310 respectively at their outlets of connection with the IP network 400.

The storage apparatus 100 comprises normally a disk drive 101, a disk adapter 103 connected to the disk drive 101 by an SCSI interface 102, a cache memory 105 connected to the disk adapter 103 by a bus 104 and a channel adapter 107 connected to the cache memory 105 by a bus 106. By the same token, the storage apparatus 200 comprises normally a disk drive 201, a disk adapter 203 connected to the disk drive 201 by an SCSI interface 202, a cache memory 205 connected to the disk adapter 203 by a bus 204 and a channel adapter 207 connected to the cache memory 205 by a bus 206. A port 108 of the channel adapter 107 is

connected to an FEC (Forward Error Correction) conversion adapter 110 by a high-speed IP interface 109. The FEC conversion adaptor 110 is connected to the IP network 400 by a high-speed IP interface 111. By the same token, a port 208 of the channel adapter 207 is connected to an FEC conversion adapter 210 by a high-speed IP interface 209. The FEC conversion adaptor 210 is connected to the IP network 400 by a high-speed IP interface 211. An example of the high-speed IP interfaces 109 and 209 is the Gigabit Ethernet (a registered trademark). The channel adapters 107 and 207 each carry out an iSCSI-protocol process. It is to be noted that the high-speed IP interface 111 connects a port of the FEC conversion adapter 110 to the IP network 400 but the port itself is not shown in the figure. By the same token, the high-speed IP interface 211 connects a port of the FEC conversion adapter 210 to the IP network 400 but the port itself is also not shown in the figure.

The host 300 includes an information-processing apparatus 301 serving as an information generation unit and an iSCSI host bus adapter (HBA) 303 connected to the information-processing apparatus 301 by an internal bus 302. The information-processing apparatus 301 has a processing unit or a memory. A port 304 of the iSCSI HBA 303 is connected to an FEC conversion adapter 310 by a high-speed IP interface 305. The FEC conversion adapter 310 is

connected to the IP network 400 by an IP interface 311.
The iSCSI HBA 303 employed in the host 300 carries out host
input/output processing and issues a command in accordance
with an instruction received from an OS (Operating System)
of the host 300.

The FEC conversion adapters 110, 210 and 310
connected to the storage apparatus 100, the storage
apparatus 200 and the host 300 respectively encode and
decode FEC codes. A detailed configuration of each of the
FEC conversion adapters 110, 210 and 310 will be described
in detail later.

It is to be noted that the IP network 400 may also
be connected to conventional storage apparatus and a
conventional host, which do not have the FEC conversion
adapters 110., 210 and 310, even though such conventional
storage apparatus and such a conventional host are not
shown in Fig. 1.

Fig. 2 is a diagram showing the internal
configuration of the FEC conversion adapter connected to a
storage apparatus in the embodiment of the present
invention. As an example of the storage apparatus, the FEC
conversion adapter 110 comprises both a transmission system
and a reception system. The transmission system includes a
physical layer 1101 on the side of the storage apparatus
100, a buffer 1102 for temporarily storing packet data

received from the storage apparatus 100, an FEC encoding unit 113 for carrying out an FEC encoding process on packet data to be transmitted, a transmission buffer 1103 for temporarily storing data completing the FEC encoding process and a physical layer 1104. The physical layer 1104 is connected to the IP network 400 by the high-speed IP interface 111.

On the other hand, the reception system comprises a physical layer 1114, a buffer 1113 for temporarily storing received data, an FEC decoding unit 1115, a buffer 1112 and a physical layer 1111. The FEC decoding unit 1115 carries out an FEC decoding process on received packet data in order to decode the data.

In addition, the FEC conversion adapter 110 also includes a control processor 116, an FEC transmission management table 117, an FEC reception management table 118 and an FEC permission table 119. The control processor 116 executes entire control in the FEC conversion adapter 110. The FEC transmission management table 117 is a table for cataloging addresses of partner apparatus to which data can be transmitted from the storage apparatus 100. Examples of such a partner are another storage apparatus and a host. The FEC permission table 119 is a table for executing management as to whether or not an FEC communication is permitted. The FEC permission table 119 is a table for

cataloging codes used for identifying a communication

partner. It is to be noted that the structures of these

tables will be described later by referring to Fig. 7.

If the address of a communication destination has

been cataloged in the FEC transmission management table 117,

data to be transmitted to the communication destination is

subjected to an FEC process in the FEC encoding unit 113

before being transferred to the IP network 400. If the

address of a communication destination has not been

cataloged in the FEC transmission management table 117, on

the other hand, data to be transmitted to the communication

destination is transferred to the IP network 400 without

being subjected to an FEC process in the FEC encoding unit

113. The FEC process of the data to be transmitted, that

is, iSCSI data, is carried out in accordance with the

communication destination's redundancy, which is also

cataloged in the FEC transmission management table 117. As

a redundancy code, it is possible to use for example a

parity-bit code produced by the commonly known XOR process

or the Reed Solomon code.

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The FEC reception management table 118 is a table

for cataloging addresses of, among others, storage

apparatus each serving as a transmission source. If the

address of a transmission source has been cataloged in the

FEC reception management table 118, the FEC decoding unit 115 carries out an FEC decoding process on data received from the transmission source in order to decode the data into iSCSI data, which is internal data. If the address of a transmission source has not been cataloged in the FEC reception management table 118, on the other hand, the data received from the transmission source is passed on to the iSCSI layer without being subjected to an FEC decoding process in the FEC decoding unit 115.

The FEC permission table 119 is a table used as a ledger for controlling communication partners. A redundancy is cataloged in the FEC permission table 119 for each iSCSI Name representing a communication partner.

Fig. 3 is a block diagram showing a storage system implemented by another embodiment of the present invention. In the case of this embodiment, in order to improve reliability, the storage apparatus 100' comprises two disk drives 101, two disk adapters 103, two cache memories 105 and two channel adapters 107. By the same token, the storage apparatus 200' comprises two disk drives 201, two disk adapters 203, two cache memories 205 and two channel adapters 207. To keep up with these double components, 2 FEC conversion adapters 110 and an FEC conversion adapter 210 are provided and connected to the IP network 400. Each of the FEC conversion adapters 110 and 210 has encoding and

decoding functions. In this embodiment, a high-speed interface 220 is used for directly connecting the port 208 of the channel adapter 207 to the IP network 400 without using an FEC conversion adapter 210 as is the case with the conventional port.

As a further embodiment, a management server 500 may be connected to the IP network 400. The management server 500 is a server for managing the FEC transmission management table 117, the FEC reception management table 118 and the FEC permission table 119, which are included in each of the FEC conversion adapters 110, 210 and 310. To put it in detail, by issuing a command to add or delete a communication destination such as a port, the management server 500 executes control as to whether or not the FEC technique is to be applied to a communication between FEC conversion adapters. The control is executed by using a management table as a table for cataloging addresses of FEC conversion adapters or ports for which communications are to be carried out by adoption of the FEC technique. It is needless to say that the addresses of a conventional storage apparatus and a conventional host are not cataloged in the management table. Thus, the effect of the FEC technique cannot be enjoyed.

Fig. 4 is a diagram showing the concept of a data transfer according to an embodiment of the present

invention. The figure shows a case in which data is
transferred from the storage apparatus 100 to the storage
apparatus 200 by way of routers A and B as well as the IP
network 400 for great distances. It is to be noted that
the destination of the data transfer is not limited to the
storage apparatus 200. Instead, a host or a server is
capable of serving as the destination of the data transfer
provided that the host or the server has an FEC conversion
adapter.

Data processed at an application layer of the
storage apparatus 100 is subjected to a protocol conversion
process at an iSCSI layer. Then, control information for
TCP and IP layers is added to the data in a process to
convert the data into a packet before the data is supplied
to the FEC conversion adapter 110 by way of a network
interface (I/F).

The FEC conversion adapter 110 carries out an FEC
process (or an encoding process) to add a redundancy code
for error correction to the data to be transmitted. This
encoding process will be described in detail later. After
the encoding process, the port number of a UDP header is
added. Then, the data is further subjected to a process to
convert the data into an IP packet before the data is
supplied to router A by way of a network I/F. Finally, the
data is output to the IP network 400 by way of route A's IP

layer and network I/F.

On the receiver side, on the other hand, the data is received and supplied to the FEC conversion adapter 210 by way of router B's IP layer and network I/F. In the FEC conversion adapter 210, the received data is supplied to an FEC layer by way of a network I/F, an IP layer and a UDP layer. At the FEC layer, the data is subjected to an FEC decoding process. This decoding process will be described later in detail. Anyway, in this FEC decoding process, the FEC redundancy code is used for correcting errors. If the number of errors is large, an ACK can by no means be returned to the transmitter side in response to the received data. For this reason, the transmitter side monitors a timeout for reception of the ACK. If the ACK is not received within a predetermined period of time, the same data is transmitted again to the receiver side. Data completing the decoding process is finally supplied to the storage apparatus 200 in which the data is supplied to an application layer by way of the same layers as the storage apparatus 100. At the application layer, the data is used in an application.

Fig. 5 is a diagram showing a typical format of the packet used in communication. On an iSCSI layer, an iSCSI PDU (Protocol Data Unit) used as a unit of data communication comprises a BHS (Basic Header Segment) and a

data segment. It is to be noted that the iSCSI PDU may include an AHS (Additional Header Sequence) between the BHS and the data segment. However, the AHS is omitted from the typical format shown in the figure. The BHS includes the length of the message. From the BHS, the start position of the data segment and the boundaries of the message are known. At the iSCSI layer, an initiator and a target carry out a communication by exchanging a message called the iSCSI PDU. The length of the iSCSI PDU is a multiple of four bytes.

At TCP, IP and data-link layers, a data link header (DLH), an IP header (IPH) and a TCP header (TCPH) are added to the head of the packet data received from the iSCSI layer described above. On the other hand, a data link trailer (DLT) is added to the tail of the iSCSI packet data. It is to be noted that, in the case of the Ethernet, an Ethernet header is used as the data link header.

The buffer (encoding buffer) 1102 shown in Fig. 2 is used for storing the headers to be sequentially added to the iSCSI packet data. The iSCSI packet data including the IP header, the TCP header and other information is subjected to a redundancy conversion process. In addition to the data link header (DLH), the IP header (IPH) and the TCP header (TCPH), a UDP header and an FEC header are further added to the data completing the redundancy

conversion process prior to a transmission.

In accordance with the present invention, in order to carry out an FEC communication, a WWN (World Wide Name), that is, an iSCSI Name, is cataloged in the FEC permission table to be used in issuing a command or the like to a communication destination. An iSCSI login is monitored. If the iSCSI login is determined to be a login for a communication destination already cataloged in the FEC permission table, an FEC communication is started by cataloging the address of the communication destination in the aforementioned FEC transmission management table and the aforementioned FEC reception management table. An IP address is typically assigned to a port on a one-to-one basis. However, a plurality of iSCSI Names can also be set for a storage apparatus. For example, a storage apparatus is divided into a plurality of partitions to be used as a plurality of disks. In this case, a plurality of iSCSI Names is set for a physical port. This way of utilization is very meaningful to the user of the storage device. A reason why communication destinations are managed by using iSCSI Names such as "iqn.1993-11.com.disk-vendor.diskarrays.sn.45678" in place of, for example, IP addresses, is that, for the user of the storage apparatus, an iSCSI Name is easier to remember than an IP address is, and it is easier to give a meaning to an iSCSI Name than an

IP address. In addition, if communication destinations are managed by using IP addresses, an IP address must be set anew in a case where a network I/F card or a network adapter card is changed or the connecting network is changed. If communication destinations are managed by using iSCSI Names, on the other hand, there is no such necessity. Moreover, in comparison with an iSCSI Name, it is quite within the bounds of possibility that an IP address is used incorrectly.

Fig. 6 is a diagram showing the format of the FEC header. The length of the FEC header is 32 bits or 4 words. Word 0 contains encoding information, the type of FEC information and an FEC control information ID. The contents of FEC control information stored in words 1, 2 and 3 change in dependence on the type of the FEC information. The encoding information is information indicating whether or not the FEC data portion has been subjected to a redundancy conversion process. The type of the FEC information is information indicating a change in information stored in the FEC permission table, a change in control table information, FEC_ACK, a report on a packet loss ratio, data and the type of redundancy data. The FEC control information ID is used for indicating control information corresponding to FEC_ACK and FEC_RJT. The FEC control information includes a data length, a redundancy

and a packet arrival ratio. The FEC data portion is included in a data packet and a redundancy packet if the FEC control packet includes an iSCSI Name.

Fig. 7 is a diagram showing typical contents of the FEC control table, the FEC transmission management table and the FEC reception management table.

The FEC permission table 119 is provided for each FEC adapter as a table functioning as the so-called ledger for executing management as to whether or not the communication partner is a target of an FEC communication. For this reason, the FEC permission table 119 is used for storing the iSCSI Name and the redundancy, which are associated with a communication partner, for each communication partner. When a login is carried out for an iSCSI Name cataloged in this FEC permission table 119, information is cataloged in the FEC transmission management table 117 or the FEC reception management table 118. It is to be noted that, for a communication partner whose iSCSI Name has not been cataloged in the FEC permission table 119, the ordinary TCP/IP communication can be carried out.

The FEC transmission management table 117 is a table used for cataloging the address, the redundancy and the encoding buffer control information for each transmission destination. The address of a transmission destination is also referred to as a destination address. In this case,

typically, an IP address is cataloged as the destination address. If a destination address has been cataloged in the FEC transmission management table 117, the iSCSI data, which is the data to be transmitted, is subjected to an FEC process carried out by the FEC encoding unit 113 and then transmitted to the IP network 400. That is to say, the data is transmitted in an FEC communication mode. If a destination address has not been cataloged in the FEC transmission management table 117, on the other hand, the iSCSI data is transmitted to the IP network 400 in a TCP/IP communication mode without being subjected to an FEC process carried out by the FEC encoding unit 113.

The FEC reception management table 118 is a table used for cataloging a packet arrival ratio and FEC decoding buffer control information for each source address. As the source address, the IP address of the transmission source is cataloged. A packet loss ratio is found from the packet arrival ratio and eventually reflected in a change in redundancy, that is, a change made by the transmission source.

It is to be noted that an iSCSI Name column can be added to each of the FEC transmission management table 117 and the FEC reception management table 118. On this column, iSCSI Names are cataloged. An iSCSI Name obtained from the FEC permission table 119 is used to search the FEC

transmission management table 117 or the FEC reception
management table 118 for the iSCSI Name's row, to or from
which information is to be added or deleted.

By referring to Figs. 8 to 10, the following
description explains operations of an FEC redundancy
conversion process, which is carried out in an FEC
conversion adapter when data is transmitted.

(A) Encoding Process

First of all, an encoding process is explained by
referring to a flowchart shown in Fig. 8. Data generated
in the storage apparatus 100 is transferred from the
channel adapter 107 to the FEC conversion adapter 110 by
way of the high-speed IP interface as a group of iSCSI
packets. At a step 801, this group of packets is received
by a physical layer 1101 of the FEC conversion adapter 110
and temporarily stored in the buffer 1102 of the FEC
conversion adapter 110. Then, at the next step 802, the
FEC encoding unit 113 checks a packet in order to form a
judgment as to whether or not the packet is either an iSCSI
login request PDU or an iSCSI login response PDU. If the
outcome of the judgment is a negation, the flow of the
process goes on to a step 803 at which the FEC transmission
management table 117 is referenced to form a judgment as to
whether or not the destination address has been cataloged
in the FEC transmission management table 117. If the

destination address has been cataloged in the FEC

transmission management table 117, the flow of the process

goes on to a step 804 to form a judgment as to whether a

storage area of the FEC encoding buffer 1103 has been

allocated. If the destination address has not been

cataloged in the FEC transmission management table 117, on

the other hand, the flow of the process goes on to a step

816 at which the packets are transferred to the physical

layer 1104 by way of the buffer 1103 as packet-group data

with the format unchanged as it is, that is, with the

TCP/IP format, without being subjected to an FEC process,

before being finally output to the IP network 400. If the

outcome of the judgment formed at the step 804 indicates

that no storage area of the FEC encoding buffer 1103 has

been allocated, the flow of the process goes on to a step

817 at which a start of a new FEC sequence is prepared.

If the FEC encoding buffer described above has been

allocated, on the other hand, the flow of the process goes

on to a step 805 to calculate an address at which the

packet is to be stored. Then, at the next step 806, the

packet is stored in the address as it is. That is to say,

the whole packet is capsulated so that the packet's IP

header, TCP header and TCP data are stored in the buffer as

they are. Subsequently, at the next step 807, an IP header,

a UDP header and an FEC header are added to the packet to

form a capsule, which is output to the IP network 400 by way of the physical layer 1104. Areas of the FEC header contain their respective pieces of information. To be more specific, an FEC header type contains FEC data and a data length contains the length of the packet, which has been capsulated. Other areas contain pieces of information according to the contents of the FEC transmission management table 117.

Then, at the next step 808, information stored in the FEC transmission management table 117 is updated. To be more specific, the FEC transmission management table 117 is updated by incrementing the number of packets for this FEC sequence by one. This updating operation is carried out to manage the total number of transmitted packets.

Subsequently, the flow of the process goes on to the next step 809 to form a judgment as to whether or not the transmitted packet is the nth packet for the FEC sequence (with a redundancy of n). If the packet is not the nth packet, the process is ended. If the packet is the nth packet, on the other hand, the flow of the process goes on to a step 810 at which redundant data is created from data stored in the FEC encoding buffer in accordance with the redundancy of n and to be finally transmitted as the (n+1)th packet. Then, at the next step 811, an IP header, a UDP header and an FEC header are added to the redundant

data and the redundant data including the IP header, the
UDP header and the FEC header is put in a capsule, which is
output to the IP network 400 by way of the physical layer
1104. Finally, at the next step 812, the sequence of FEC
processing operations is ended. If the FEC retransmission
function is provided, the sequence of FEC processing
operations is ended by waiting for an FEC_ACK. If the FEC
retransmission function is not provided, on the other hand,
the sequence of FEC processing operations is ended by
making the FEC encoding buffer available.

If the outcome of the judgment formed at the step
802 is an affirmation, on the other hand, the flow of the
process goes on to a step 813 to form a judgment as to
whether or not the group of packets is an iSCSI initial
login request PDU. If the outcome of the judgment is an
affirmation, the flow of the process goes on to a step 814
at which the FEC permission table is checked in order to
form a judgment as to whether or not the target name has
been cataloged in the table. If the target name has been
cataloged in the FEC permission table, the flow of the
process goes on to a step 815 at which the destination
address of the packet is cataloged in the FEC transmission
management table before proceeding to the next processing.
If the outcome of the judgment formed at the step 813 is a
negation, on the other hand, the flow of the process goes

on to a step 816 at which the packet is transferred to the physical layer 1104 as it is without being subjected to the FEC processing. Finally, the process is ended.

(B) Processing Operations for Insufficient Data to Be Transmitted

By referring to a flowchart shown in Fig. 9, the following description explains processing operations, which are carried out in a case where data to be transmitted is insufficient. This process is carried out to form a judgment as to whether or not the normal number of packets has been received from the channel adapter.

When a transmission of a packet group in a certain transmission sequence is started, a timer interrupt is generated in order to commence a counting operation of a timer at a step 820. Then, at the next step 821, the timer is checked in order to form a judgment as to whether or not a predetermined period of time has lapsed since the start of the FEC sequence process to transmit data. If the predetermined period of time has not lapsed, the process is ended. If the predetermined period of time has lapsed, on the other hand, the flow of the process goes on to a step 822 at which redundant data with an FEC sequence redundancy of m is created for m packets already received from the physical layer 1101 in this FEC sequence. It is to be noted that the redundancy of m is stored in the FEC header

of the redundant packet.

(C) Operations to Retransmit FEC Packets

Next, operations to retransmit an FEC packet are explained by referring to a flowchart shown in Fig. 10. In the basic sequence of this process, when an FEC packet is received on the receiver side, an FEC_ACK is transmitted to the transmitter side in response to the packet. By receiving this FEC_ACK, the transmitter side is capable of recognizing that the transmitted packet has been received correctly by the partner.

When the transmitter side transmits an FEC packet, timer interrupt processing is carried out at a step 830. Then, at the next step 831, the timer is checked in order to form a judgment on a timeout for the FEC sequence in a state of waiting for an FEC_ACK. If the outcome of the judgment indicates that a timeout has not occurred, the process is ended. If the outcome of the judgment formed at the step 831 indicates that a timeout has occurred, on the other hand, the flow of the process goes on to a step 832 at which processing is carried out to retransmit all packets for this FEC sequence. By the same token, the timer is reset to start the operation to measure the time to wait for an FEC_ACK for all the retransmitted packets.

By referring to a flowchart shown in Figs. 11 and 12, the following description explains operations of an FEC

decoding process, which is carried out by the FEC

conversion adapter 110 when data is received. The

flowchart shown in Figs. 11 and 12 begins with a step 901

at which a packet is received from the physical layer 1114.

Then, at the next step 902, the packet is checked in order

to form a judgment as to whether or not the packet is

either an iSCSI login request PDU or an iSCSI login

response PDU. If the outcome of the judgment is an

affirmation, the flow of the process goes on to a step 919

to form a judgment as to whether or not this packet is an

iSCSI initial login request. If the outcome of the

judgment is an affirmation, the flow of the process goes on

to a step 920 to form a judgment as to whether or not the

source name (that is, the iSCSI Name) has been cataloged in

the FEC permission table. If the source name has been

cataloged in the FEC permission table, the flow of the

process goes on to a step 921 at which the destination

address of the packet is cataloged in the FEC reception

management table.

If the outcome of the judgment formed at the step

902 is a negation, on the other hand, the flow of the

process goes on to a step 903 to form a judgment as to

whether or not the destination port number has been

provided for FEC communication use. If the outcome of the

judgment formed at the step 903 is a negation, the flow of

the process goes on to a step 922 at which the packet is transferred to the physical layer 1111 as it is without being subjected to an FEC process. If the outcome of the judgment formed at the step 903 indicates that the destination port number has been provided for FEC communication use, on the other hand, the flow of the process goes on to a step 904 to form a judgment as to whether or not the packet is a packet for changing information cataloged in the FEC permission table. If the outcome of the judgment is an affirmation, the flow of the process goes on to a step 923 at which the contents of the FEC permission table are updated. Then, at the next step 924, TCP processing is carried out. That is to say, an ACK is transmitted as a response and the process is ended.

If the outcome of the judgment indicates that the packet is not a TCP packet, on the other hand, the flow of the process goes on to a step 905 at which the packet is checked in order to form a judgment as to whether or not the packet is an UDP packet. If the packet is an UDP packet, the flow of the process goes on to a step 906 to form a judgment as to whether or not the source address of the packet has been cataloged in the FEC reception management table. If the source address of the packet has been cataloged in the FEC reception management table, the flow of the process goes on to a step 907 to form a

judgment as to whether or not the packet has been destroyed.
If the packet has not been destroyed, the flow of the
process goes on to a step 908 to form a judgment as to
whether or not the packet is a packet of an FEC sequence
for which a timeout has occurred. If the outcome of the
judgment is an affirmation, the flow of the process goes on
to a step 925 at which the packet arrival ratio for the
packet is found from the FEC reception management table by
using the source address and updated. Then, at the next
step 926, the packet is discarded and the process is ended.

If the outcome of the judgment formed at the step
908 indicates that the packet is not a packet of an FEC
sequence for which a timeout has occurred, on the other
hand, the flow of the process goes on to a step 909 at
which the packet arrival ratio for the packet is found from
the FEC reception management table by using the source
address and updated. Then, the flow of the process goes on
to a step 910 at which the packet is checked in order to
form a judgment as to whether or not the packet includes an
ID indicating a decoded FEC sequence. If the outcome of
the judgment is an affirmation, the flow of the process
goes on to the step 926 at which the packet is discarded
and the process is ended. If the packet does not include
an ID indicating a decoded FEC sequence, on the other hand,
the flow of the process goes on to a step 911 at which the

FEC decoding buffer is checked in order to form a judgment as to whether or not an area of the FEC decoding buffer has been allocated to the FEC sequence. If the outcome of the judgment is a negation, on the other hand, the flow of the process goes on to the step 927 at which processing to decode the FEC sequence is carried out. In this processing, typically, an available area in the FEC decoding buffer is allocated and a timer is set.

If the outcome of the judgment formed at the step 911 is an affirmation, the flow of the process goes on to a step 912 to compute an address at which the packet is to be stored. Then, at the next step 913, the packet is stored at the address. Subsequently, at the next step 914, the packet is checked in order to form a judgment as to whether or not the packet is an FEC control packet. If the outcome of the judgment is a negation, the flow of the process goes on to a step 928 at which the packet is checked in order to form a judgment as to whether or not the packet is a data packet. If the packet is a data packet, the flow of the process goes on to a step 929 at which the packet is transferred to the physical layer 1111 as a packet with its data capsulated.

Then, the flow of the process goes on to a step 915 to form a judgment as to whether or not the packet is received to the nth in the FEC sequence of the packet,

where the redundancy of the FEC is n in this case.   If
the outcome of the judgment is an affirmation, the flow of
the process goes on to a step 916 to form a judgment as to
whether or not the FEC sequence's packet not received is a
redundant packet. If the packet not received is a
redundant packet, the flow of the process goes on to a step
917 at which an FEC_ACK is transmitted in response to the
packet and a sequence ID is stored. Then, at the next step
918, the process to handle this received FEC sequence is
ended. To put it in detail, operations are carried out to
de-allocate the allocated area of the FEC decoding buffer
and temporarily store the sequence ID in the FEC reception
management table as an ID of a processed sequence.

If the outcome of the judgment formed at the step
916 to determine whether or not the FEC sequence's packet
not received is a redundant packet is not an affirmation,
on the other hand, the flow of the process goes on to a
step 930 at which the packet not received is reconstructed
in accordance with a redundancy of n. It is to be noted
that, if the FEC header of the redundant packet generated
at the step 822 shows a redundancy of m, the packet not
received is also reconstructed in accordance with the
redundancy of m too. Then, the flow of the process goes on
to the next step 931 to form a judgment as to whether or
not the reconstructed packet is an FEC control packet. If

the outcome of the judgment is not an affirmation, the flow of the process goes on to a step 932 to form a judgment as to whether or not the reconstructed packet is an FEC data packet. If the reconstructed packet is an FEC data packet, the flow of the process goes on to a step 934 at which a packet with the reconstructed data capsulated is transferred to the physical layer 1111. It is to be noted that processing of an FEC control packet will be described later.

By referring to flowcharts shown in Figs. 13 to 15, the following description explains operations of a process carried out in the event of an error detected in reception of data and a process related to loss of a packet.

(A) Processing of Insufficient Received Data

This processing is a process, which is carried out when data of an FEC sequence completing an FEC redundancy process is not sufficient so that the FEC sequence cannot be decoded within a predetermined period of time.

The flowchart shown in Fig. 13 begins with a step 1501 at which a timer interrupt is generated when reception of a packet of a certain sequence is started. Then, the flow of the process goes on to the next step 1502 to form a judgment as to whether or not a predetermined period of time has lapsed since the start of the process before the nth packet of a certain sequence with a redundancy of n is

received. If the period of time has lapsed before the nth

packet is received, the flow of the process goes on to a

step 1503 to compute a variable x representing the number

of FEC-sequence packets not received. Typically, the

variable x is computed by subtracting the number of

received packets from n (n - the number of received

packets). Subsequently, at the next step 1504, the

decoding buffer used for storing the FEC sequence is

cleared. Then, at the next step 1505, the packet arrival

ratio is computed and the process is ended. That is to say,

the FEC sequence's packets not received are regarded as

lost packets in computing the packet arrival ratio. A

packet arrival ratio obtained as a result of the

computation is used as an updated packet arrival ratio.

(B) Process to Report Packet Arrival Ratio

The computed packet arrival ratio is cataloged in

the FEC reception management table on the receiver side and

also transmitted to the transmitter side as well. On the

transmitter side, a packet loss ratio is used to find a

redundancy, which is cataloged in the FEC transmission

management table.

The flowchart shown in Fig. 14 begins with a step

1601 at which a timer interrupt is generated on the

receiver side. Then, at the next step 1602, a management

frame used for reporting a packet arrival ratio is

transmitted for each of FEC source addresses. Subsequently,
at the next step 1603, a packet used for reporting the
packet arrival ratio is stored in the FEC encoding buffer.
The packet arrival ratio (Num_Alive PKT) is a ratio of the
number of FEC data packet received prior to the occurrence
of a timeout without being destroyed to the number of
transmitted FEC data packets.

 (C): Processing of FEC Control Packets

        The flowchart shown in Fig. 15 begins with a step
1701 at which a received packet is checked in order to form
a judgment as to whether or not the packet is a packet used
for reporting a packet loss ratio. If the received packet
is a packet used for reporting a packet loss ratio, the
flow of the process goes on to a step 1702 at which a
redundancy for a communication destination is computed from
the reported packet loss ratio and a safety margin
determined in advance. The computed redundancy is
cataloged in the FEC transmission management table as an
update value of a corresponding entry. If the outcome of
the judgment formed at the step 1701 indicates that the
received packet is not a packet used for reporting a packet
loss ratio, on the other hand, the flow of the process goes
on to a step 1704 at which the received packet is checked
in order to form a judgment as to whether or not the packet
is a packet used for changing information stored in the FEC

transmission management table. If the outcome of the

judgment formed at the step 1704 is an affirmation, the

flow of the process goes on to a step 1705 at which the

information stored in the FEC transmission management table

is updated. If the received packet is not a packet used

for changing information stored in the FEC transmission

management table, on the other hand, the flow of the

process goes on to a step 1706 at which the received packet

is checked in order to form a judgment as to whether or not

the packet is a packet containing FEC_ACK. If the outcome

of the judgment is an affirmation, the flow of the process

goes on to a step 1707 at which the processing of the

transmitted FEC sequence is ended. In this processing, the

allocated area of the FEC encoding buffer is released, the

state of waiting for FEC_ACK is terminated and the relevant

information is deleted from the FEC transmission management

table.

Embodiments of the present invention have been

explained so far. However, the scope of the present

invention is not limited to the embodiments. That is to

say, a variety of changes can be made to the embodiments.

Fig. 16 is a block diagram showing an embodiment

implementing a modified configuration of the channel

adapter employed in the storage apparatus. This channel

adapter 107 also includes some of the functions of the FEC

conversion adapter described earlier. To be more specific,
the channel adapter 107 also includes components of the FEC
conversion adapter, except components on the channel-
adapter side, in addition to an SCSI command control unit
1071, a protocol control unit 1072, a transmission buffer
1073 and a reception buffer 1074. The transmission buffer
1073 and the reception buffer 1074 are employed as common
buffers. The excluded components on the channel-adapter
side are the buffers 1102 and 1112 and the physical layers
1101 and 1111. The remaining configuration is the same as
that shown in Fig. 2.

In addition, in the embodiments described above,
data is transmitted by changing the redundancy to a value
suitable for a transmission destination in accordance with
the packet loss ratio. However, it is also possible to
change intervals, at which packets are transmitted, in
accordance with the packet loss ratio. In a broad sense,
the data transmission state is changed in accordance with
the packet loss ratio.

In accordance with the preferred embodiments, it is
possible to provide a storage apparatus conforming to an
iSCSI protocol adopting the FEC technique. In addition, by
feeding a packet loss state grasped on the receiver side
back to the transmitter side, the transmitter side is
capable of transmitting data by changing the state of FEC

processing and the data transmission redundancy to a value
suitable for a transmission destination whereas the
receiver side is capable of decoding the received data
accordingly. As a result, it is possible to provide a
storage system capable of restoring data even if packets
have been lost.